

# A 2-categorical Quillen's Theorem A

Part 1

(Work by  
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Higher Category Reading Seminar, nty  
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## Plan

- Idea and the motivation
- Basic 2-cat thg
- Intro
- Dwyer Nerve
- lax slices and functors
- ad-weakness

## Thm (Quillen Thm A)

$C, D \in \text{Cat}$  and  $F: C \rightarrow D$ .

$\forall$  object  $n$  of  $D$  the full subcategory  $|N(F_n)|$  of the comma cat.  $F/n$  is contractible then

$|N(C)| \rightarrow |N(D)|$  is a weak hty eq.

## A version of Quillen's Thm A

Thm  $X, Y \in \text{cat}$ ,  $F: X \rightarrow Y$  is a functor

Then if  $\forall n \in Y$

$\exists n' \in X$  and a map  $n \rightarrow F(n')$

and every morphism  $n \rightarrow F(z)$  represent an initial obj of  $|N(F_n)|$

Then  $|F|: |N(C)| \rightarrow |N(D)|$  is a hty eq.

## 2-cat thg

Abuse of notation  $\mathbb{Z}\text{-cat}$  means a strict 2-cat.

Note  $X \in \mathbb{Z}\text{cat}$

1-morph	is denoted	$X(0p, -)$
2-morph	"	$X(-, \alpha)$
1,2-morph	"	$X(0p, \alpha)$

Def.  $X, Y \in \text{Cat}$ . A normal lax 2-functor  $F: X \rightarrow Y$

- a map  $F: \text{ob}(X) \rightarrow \text{ob}(Y)$

-  $\forall$  pairs  $x, x' \in X$ , a functor

$$F_{x, x'} : X(x, x') \rightarrow X(F(x), F(x'))$$

-  $\forall$  triple  $x, x', x'' \in X$

$$X(x, x'') \times X(x, x') \xrightarrow{\circ} X(x, x')$$

$$\begin{array}{ccc} & \xleftarrow{G} & \\ F \downarrow & & \downarrow F \end{array}$$

$$Y(F(x), F(x'')) \times Y(F(x), F(x')) \xrightarrow{\circ} Y(F(x), F(x'))$$

$\exists$  a nat'l trans.  $G$  (compositor)

st.

$$F_{x, x'}(Id_x) = Id_{F(x)} \quad \forall x \in X$$

- the compositor satisfies the hexagon identity

$$G_{x, Id_x} = Id_{F(x)} \quad \Delta \quad G_{Id_x, x} = Id_{F(x)}$$

Lcat  $\rightarrow$  cat. of 2-cats with normal lax 2-fun.

Def. We get a cosimp. obj. in Lcat by composing

$$j^0: \Delta \rightarrow \text{cat} \quad \text{with } j: \text{cat} \rightarrow \text{Lcat}$$

(Theorem of math: we don't check  $j \circ j^0$  by  $j^0$ )

Then we get a functor with this cosimp. obj.

$$N_2: \text{Lcat} \rightarrow \text{Set} \quad \text{with } N_2(x) =$$

$$= \text{Lcat}(F(x), x)$$

(Dustin Nerve)

Def. Let  $I$  be a linearly ord. fin. set, then let's define  $\Theta^I$ .

obj: elements of  $I$

- The cat  $\Theta^I(x, y)$  of morph.  $x, y \in I$  is defined as a poset of fin. sub  $S \subseteq I$  all  $\min(S) = x, \max(S) = y$

- The comp. is given by  $x, y, z \in I$

$$\Theta^I(x, y) \times \Theta^I(y, z) \rightarrow \Theta^I(x, z)$$

$$\text{(Notice } \Theta^{[n]} \cong \Theta^1)$$

(as a poset) obj  $\emptyset$

The compositor  $\Delta \xrightarrow{\Theta^1} \text{Zcat} \rightarrow \text{Lcat}$  is  $\cong$   $\Theta^1$  should be

$\forall n \in \mathbb{N}$ , we can define

$$h_n: [n] \rightarrow \Theta^1 \quad \text{def on } [n-1]$$

$$h_n(x, y) = \{x, y\}$$

Prop.  $X \in \text{Zcat}$

(1)  $N_2$  is fully faithful

(2) The lax functor  $N_2$  forms a nat'l trans of cosimp. obj.

(3)  $\forall n \in \mathbb{N}$  the lax functor

$h_n$  induces an isoch.

$$\text{Lcat}([n], Y) \cong \text{Zcat}(\Theta^1, Y)$$

# Lax slices and functors

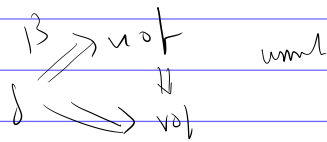
Def.  $X \in \mathcal{Z}\text{Cat} \rightarrow \alpha \in X$

Let's define  $X_{\alpha} \subset \mathcal{Z}\text{Cat}$ , the lax slice cat, object.

obj of  $X_{\alpha}$  are morph  $f: x \rightarrow y$  in  $X$

Morph are from  $f: x \rightarrow y$   $g: v \rightarrow z$  ~~are omitted~~  
 a morph  $h: v \rightarrow x$  in  $X$  with a 2-morph  
 $\beta: g \Rightarrow \alpha \circ f$ .

a 2-morph  $\eta$  from  $(\alpha, \beta)$ ,  $(\alpha, \alpha)$  unital 2-morph  
 $\eta: \alpha \Rightarrow \alpha$



We can define the of lax slice cat  $X_{\alpha}$  as

$$\left( (X^{(-, \alpha)})_{\alpha} \right)^{(-, \alpha)}$$

Prop:  $\mathcal{L}\text{Cat} \rightarrow \text{Cat}$  is a functor  
 $(X, \alpha) \mapsto X_{\alpha}$

## $\infty$ -braid

Def. A marked 2-cat is a pair  $X^{\dagger} = (X, W_X)$  where  
 $X \in \mathcal{Z}\text{Cat} \rightarrow W_X \subset \text{Mor}(X)$  marking all idemp.

A functor of marked 2-cats  $F: X^{\dagger} \rightarrow Y^{\dagger}$  is a 2-functor  
 $F: X \rightarrow Y$  st.  $F(W_X) \subset W_Y$

Notat.  $\mathcal{Z}\text{Cat}^{\dagger} = \text{cat of marked 2-cats}$

Def. A cat with weps is a pair  $X^{\dagger} = (X, W_X)$  where  $X \in \text{Cat}$   
 $\rightarrow W_X \subset X$  is a wide subcat.

A homotopy of functors  $F: X^{\dagger} \rightarrow Y^{\dagger}$  is a pair  
 $(F: X \rightarrow Y)$  st. it maps  $W_X$  to  $W_Y$

Analogously, a 2-cat with weps  $X^{\dagger}$  consists of a 2-cat  
 $X$  with the str. of a cat with weps on the  
 underlying 1-cat.

Simply we define a homotopy 2-functor.

Notat.  $\mathcal{Z}\text{Cat}^{\text{wep}} = \text{cat of 2-cats with weps and homotopy 2-functors}$

Def. A marked 2-cat  $(X, W_X)$  is called ret-well if

1)  $W_X$  contains all eqs of  $X$

2)  $(X, W_X)$  is a 2-cat with eqs.

3) For  $f \in W_X, g \in X$  if we're an invertible 2-mor  $f \xrightarrow{\cong} g$  then  $f \in W_X$

Def. For a formula  $f: X^+ \rightarrow Y^+$  of marked 2-cats,  $\forall y \in Y^+$ , we can define a marked 2-cat  $X_y^+$  where  $X_y^+$  is the best slice cat with an edge marked iff the assoc. 2-mor is invertible if the assoc. 2-mor is marked in  $X^+$ .

Def.  $(X, W) \in \text{Set}^+$ . An  $\omega$ -cat localization of  $X$  by  $W$  is an  $\omega$ -cat  $L_W(X)$  eq(1) with map  $\mathcal{D}_X: X \rightarrow L_W(X)$  in  $\text{Set}^+$ , st.  $\forall e \in \overline{\omega\text{-cat}}$  the map  $\mathcal{D}_X^*: \text{Fun}(X, W), \mathcal{L} \xrightarrow{\cong} \text{Fun}(L_W(X), \mathcal{L})$  is an eq of  $\omega$ -cat.

Ex 7.1.10  $\text{Fun}(X, W) \in \text{Set}^+$ , the best.

Ex.  $X \rightarrow L_W(X)$  is adjoint & universal.

factorial  $f: X \rightarrow Y$  st. preserve with / preserve  $\omega$ -cat

Thanks you